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# Efficiency of Embedded Explicit Pseudo Two-Step RKN Methods on a Shared Memory Parallel Computer*

**N. H. Cong[1], H. Podhaisky[2], and R. Weiner[2]**

[1]*Faculty of Math., Mech. and Inform., Hanoi University of Science*
*334 Nguyen Trai, Thanh Xuan, Hanoi, Vietnam*
[2]*FB Mathematik und Informatik, Martin-Luther-Universität Halle-Wittenberg*
*Theodor-Lieser-Str. 5, D-06120 Halle, Germany*

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**Abstract.** The aim of this paper is to construct two embedded explicit pseudo two-step RKN methods (embedded EPTRKN methods) of order 6 and 10 for nonstiff initial-value problems (IVPs) $\mathbf{y}''(t) = \mathbf{f}(t, \mathbf{y}(t))$, $\mathbf{y}(t_0) = \mathbf{y}_0$, $\mathbf{y}'(t_0) = \mathbf{y}'_0$ and investigate their efficiency on parallel computers. For these two embedded EPTRKN methods and for expensive problems, the parallel implementation on a shared memory parallel computer gives a good speed-up with respect to the sequential one. Furthermore, for numerical comparisons, we solve three test problems taken from the literature by the embedded EPTRKN methods and the efficient nonstiff code ODEX2 running on the same shared memory parallel computer. Comparing computing times for accuracies received shows that the two new embedded EPTRKN methods are superior to the code ODEX2 for all the test problems.

## 1. Introduction

The arrival of parallel computers influences the development of numerical methods for a nonstiff initial-value problem (IVP) for systems of special second-order ordinary differential equations (ODEs)

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$$\mathbf{y}''(t) = \mathbf{f}(t, \mathbf{y}(t)), \quad \mathbf{y}(t_0) = \mathbf{y}_0, \quad \mathbf{y}'(t_0) = \mathbf{y}'_0, \qquad \mathbf{y}, \mathbf{f} \in \mathbb{R}^d. \tag{1.1}$$

The most efficient numerical methods for solving this problem are the explicit Runge-Kutta-Nyström (RKN) and extrapolation methods. In the literature, sequential explicit RKN methods up to order 11 can be found in e.g., [16-21, 23, 28]. In order to exploit the facilities of parallel computers, a number of parallel explicit methods have been investigated, for example in [2-6, 9-14]. A common challenge in the latter mentioned works is to reduce, for a given order of accuracy, the required number of effective sequential **f**-evaluations per step, using parallel processors.

In previous work of Cong et al. [14], a general class of explicit pseudo two-step RKN methods (EPTRKN methods) for solving problems of the form (1.1) has been investigated. These EPTRKN methods are ones of the cheapest parallel explicit methods in terms of number of effective sequential **f**-evaluations per step. They can be easily equipped with embedded formulas for a variable stepsize implementation (cf. [9]). With respect to the number of effective sequential **f**-evaluations for a given accuracy, the EPTRKN methods have been shown to be much more efficient than most efficient sequential and parallel methods currently available for solving (1.1) (cf. [9, 14]).

Most numerical comparisons of parallel and sequential methods are done by means of the number of effective sequential **f**-evaluations for a given accuracy on a sequential computer ignoring the communication time between processors (cf. e.g., 1, 3, 5, 6]). In comparisons of different codes running on parallel computers, the parallel codes often give disappointing results. However, in our recent work [15], two parallel codes EPTRK5 and EPTRK8 of oder 5 and 8, respectively, have been proposed. These codes are based on EPTRK methods considered in [7, 8] which are a "first-order" version of the EPTRKN methods. The EPTRK5 and EPTRK8 codes have been shown to be more efficient than the codes DOPRI5 and DOP853 for solving expensive nonstiff first-order problems on a shared memory parallel computer. We have also obtained a similar performance of a parallel implementation of the BPIRKN codes for nonstiff special second-order problems (see [13]). These promising results encourage us to pursue the efficiency investigation of a real implementation of the EPTRKN methods on a parallel computer. This investigation consists of choosing relatively good embedded EPTRKN methods, defining reasonable error estimate for stepsize strategy and comparing the resulting EPTRKN methods with the code ODEX2 which is among the most efficient sequential nonstiff integrators for special second-order ODE systems of the form (1.1). Differed from the EPTRKN methods considered in [9], the embedded EPTRKN methods constructed in this paper are based on collocation vectors which minimise the stage error coefficients and/or satisfy the orthogonality relation (see Sec. 3.1). In addition to that, their embedded formulas are also derived by a new way (see Sec. 2.2). Although the class of EPTRKN methods contains methods of arbitrary high order, we consider only two EPTRKN methods of order 6 and 10 for numerical comparisons with the code ODEX2.

We have to note that the choice of an implementation on a shared memory

parallel computer is due to the fact that such a computer can consist of several processors sharing a common memory with fast data access requiring less communication times, which is suited to the features of the EPTRKN methods. In addition, there are the advantages of compilers which attempt to parallelize codes automatically by reordering loops and sophisticated scientific libraries (cf. e.g., [1]).

In order to see a possible speed-up of a parallel code, the test problems used in Sec. 3 should be expensive. Therefore, the relatively small problems have been enlarged by scaling.

## 2. Variable Stepsize Embedded EPTRKN Methods

The EPTRKN methods have been recently introduced and investigated in [9, 14]. For an implementation with stepsize control, we consider variable stepsize embedded EPTRKN methods. Because EPTRKN methods are of a two-step nature, there is an additional difficulty in using these methods with variable stepsize mode. We overcome this difficulty by deriving methods with variable parameters (cf. e.g., [24, p. 397; 1, p. 44]). Thus, we consider the variable stepsize EPTRKN method (cf. [9])

$$\mathbf{Y}_n = \mathbf{e}\otimes \mathbf{y}_n + h_n\mathbf{c}\otimes \mathbf{y}'_n + h_n^2(A_n\otimes I)\mathbf{F}(t_{n-1}\mathbf{e}+h_{n-1}\mathbf{c}, \mathbf{Y}_{n-1}), \tag{2.1a}$$

$$\mathbf{y}_{n+1} = \mathbf{y}_n + h_n\mathbf{y}'_n + h_n^2(\mathbf{b}^T\otimes I)\mathbf{F}(t_n\mathbf{e}+h_n\mathbf{c}, \mathbf{Y}_n),$$

$$\mathbf{y}'_{n+1} = \mathbf{y}'_n + h_n(\mathbf{d}^T\otimes I)\mathbf{F}(t_n\mathbf{e}+h_n\mathbf{c}, \mathbf{Y}_n), \tag{2.1b}$$

with variable stepsize $h_n = t_{n+1} - t_n$ and variable parameter matrix $A_n$. This EPTRKN method is conveniently specified by the following tableau:

| $A_n$ | $\mathbf{c}$ | $O$ |
|---|---|---|
| $\mathbf{0}^T$ | $\mathbf{y}_{n+1}$ | $\mathbf{b}^T$ |
| $\mathbf{0}^T$ | $\mathbf{y}'_{n+1}$ | $\mathbf{d}^T$ |

At each step, $2s$ **f**-evaluations of the components of the big vectors $\mathbf{F}(t_{n-1}\mathbf{e} + h_{n-1}\mathbf{c}, \mathbf{Y}_{n-1}) = (\mathbf{f}(t_{n-1} + c_i h_{n-1}, \mathbf{Y}_{n-1,i}))$ and $\mathbf{F}(t_n\mathbf{e} + h_n\mathbf{c}, \mathbf{Y}_n) = (\mathbf{f}(t_n + c_i h_n, \mathbf{Y}_{n,i}))$, $i = 1,\ldots,s$ are used in the method. However, $s$ **f**-evaluations of the components of $\mathbf{F}(t_{n-1}\mathbf{e} + h_{n-1}\mathbf{c}, \mathbf{Y}_{n-1})$ are already available from the preceding step. Hence, we need to compute only $s$ **f**-evaluations of the components of $\mathbf{F}(t_n\mathbf{e} + h_n\mathbf{c}, \mathbf{Y}_n)$, which can be done in parallel. Consequently, on an $s$-processor computer, just one **f**-evaluation is required per step. In this way, parallelization in an EPTRKN method is achieved by sharing the **f**-evaluations of $s$ components of the big vector $\mathbf{F}(t_n\mathbf{e} + h_n\mathbf{c}, \mathbf{Y}_n)$ over a number of available processors. An additional computational effort consists of a recomputation of the variable parameter matrix $A_n$ defined by (2.2e) below when the stepsize is changed.

### 2.1. Method Parameters

The matrix $A_n$ and the weight vectors $\mathbf{b}^T$ and $\mathbf{d}^T$ of the method (2.1) are derived by the order conditions (see [9, 14])

$$\frac{\tau_n^{j-1}\mathbf{c}^{j+1}}{j+1} - A_n j(\mathbf{c}-\mathbf{e})^{j+1} = \mathbf{0}, \quad j = 1,\dots,q, \tag{2.2a}$$

$$\frac{1}{j+1} - \mathbf{b}^T j\mathbf{c}^{j-1} = 0, \quad j = 1,\dots,p, \tag{2.2b}$$

$$\frac{1}{j} - \mathbf{d}^T\mathbf{c}^{j-1} = 0, \quad j = 1,\dots,p, \tag{2.2c}$$

where $\tau_n = h_n/h_{n-1}$ is the stepsize ratio. Notice that the conditions (2.2b), (2.2c) for $p = s$ define the weight vectors of a direct collocation-based IRKN method (cf. [26]).

For $q = p = s$, by defining the matrices and vectors

$$P = \Big(\frac{c_i^{j+1}}{j+1}\Big), \quad Q = \big(j(c_i-1)^{j-1}\big), \quad R = \big(jc_i^{j-1}\big), \quad S = \big(c_i^{j-1}\big),$$
$$D_n = \mathrm{diag}(1,\tau_n,\dots,\tau_n^{s-1}), \quad \mathbf{v} = \Big(\frac{1}{j}\Big), \quad \mathbf{w} = \Big(\frac{1}{j+1}\Big), \quad i,j = 1,\dots,s,$$

the conditions (2.2) can be written in the form

$$A_nQ - PD_n = O, \quad \mathbf{w}^T - \mathbf{b}^TR = \mathbf{0}^T, \quad \mathbf{v}^T - \mathbf{d}^TS = \mathbf{0}^T, \tag{2.2d}$$

which implies the explicit formulas for the parameters of a EPTRKN method

$$A_n = PD_nQ^{-1}, \quad \mathbf{b}^T = \mathbf{w}^TR^{-1}, \quad \mathbf{d}^T = \mathbf{v}^TS^{-1}. \tag{2.2e}$$

For determining the order of EPTRKN methods constructed in Sec. 3.1, we need the following theorem, which is similar to Theorem 2.1 in [9].

**Theorem 2.1.** *If the step ratio $\tau_n$ is bounded from above (i.e., $\tau_n \leqslant \Omega$) and if function $\mathbf{f}$ is Lipschitz continuous, the s-stage EPTRKN method (2.1) with parameter matrix and vectors $A_n$, $\mathbf{b}$, $\mathbf{d}$ defined by (2.2e) is of stage order $q = s$ and order $p = s$ for any collocation vector $\mathbf{c}$ with distinct abscissae $c_i$. It has higher stage order $q = s+1$ and order $p = s+1$ or $p = s+2$ if in addition the orthogonality relation*

$$P_j(1) = 0, \quad P_j(x) := \int_0^x \xi^{j-1}\prod_{i=1}^{s}(\xi - c_i)d\xi, \quad j = 1,\dots,k,$$

*holds for $k = 1$ or $k \geqslant 2$, respectively.*

The proof of this theorem follows the same line as in the proof of a very similar theorem formulated in [9, proof of Theorem 2.1].

2.2. Embedded Formulas

With the aim to have a cheap error estimate used in the stepsize selection for an implementation of EPTRKN methods with stepsize control, we shall equip the $p$th-order EPTRKN method (2.1) with the following embedded formula

$$\begin{aligned}\widehat{\mathbf{y}}_{n+1} =& \mathbf{y}_n + h_n\mathbf{y}'_n + h_n^2(\widehat{\mathbf{b}}^T \otimes I)\mathbf{F}(t_n\mathbf{e} + h_n\mathbf{c}, \mathbf{Y}_n),\\ \widehat{\mathbf{y}}'_{n+1} =& \mathbf{y}'_n + h_n(\widehat{\mathbf{d}}^T \otimes I)\mathbf{F}(t_n\mathbf{e} + h_n\mathbf{c}, \mathbf{Y}_n),\end{aligned} \tag{2.3}$$

where, the weight vectors $\widehat{\mathbf{b}}$ and $\widehat{\mathbf{d}}$ are determined by satisfying the following conditions which come from (2.2b) and (2.2c)

$$\frac{1}{j+1} - \widehat{\mathbf{b}}^T j\mathbf{c}^{j-1} = 0, \quad j = 1, \ldots, s-2, \quad \frac{1}{s} - \widehat{\mathbf{b}}^T(s-1)\mathbf{c}^{s-2} \neq 0, \tag{2.4a}$$

$$\frac{1}{j} - \widehat{\mathbf{d}}^T\mathbf{c}^{j-1} = 0, \quad j = 1, \ldots, s-1, \quad \frac{1}{s} - \widehat{\mathbf{d}}^T\mathbf{c}^{s-1} \neq 0. \tag{2.4b}$$

In the two EPTRKN codes considered in this paper, we use these embedded weight vectors defined as

$$\widehat{\mathbf{b}}^T = \Big(\mathbf{w}^T - \frac{1}{10}\mathbf{e}_{s-1}^T\Big)R^{-1}, \quad \widehat{\mathbf{d}}^T = \Big(\mathbf{v}^T - \frac{1}{10}\mathbf{e}_s^T\Big)S^{-1}, \tag{2.5}$$

where $\mathbf{e}_s^T = (0, \ldots, 0, 1)$ and $\mathbf{e}_{s-1}^T = (0, \ldots, 1, 0)$ are the $s$-th and $(s-1)$-th unit vectors. It can be seen that the following simple theorem holds

**Theorem 2.2.** *The embedded formula defined by* (2.3) *and* (2.5) *is of order* $s-1$ *for any collocation vector* $\mathbf{c}$ *with distinct abscissae* $c_i$.

In this way we have an estimate for the local error of order $\widehat{p} = s-1$ without additional $\mathbf{f}$-evaluations given by

$$\mathbf{y}_{n+1} - \widehat{\mathbf{y}}_{n+1} = O(h_n^{\widehat{p}+1}), \quad \mathbf{y}'_{n+1} - \widehat{\mathbf{y}}'_{n+1} = O(h_n^{\widehat{p}+1}). \tag{2.6}$$

Thus, we have defined the embedded EPTRKN method of orders $p(\widehat{p})$ given by (2.1), (2.2e), (2.3) and (2.5) which can be specified by the tableau

| $A_n$ | $\mathbf{c}$ | $O$ |
|---|---|---|
| $\mathbf{0}^T$ | $\mathbf{y}_{n+1}$ | $\mathbf{b}^T$ |
| $\mathbf{0}^T$ | $\mathbf{y}'_{n+1}$ | $\mathbf{d}^T$ |
| $\mathbf{0}^T$ | $\widehat{\mathbf{y}}_{n+1}$ | $\widehat{\mathbf{b}}^T$ |
| $\mathbf{0}^T$ | $\widehat{\mathbf{y}}'_{n+1}$ | $\widehat{\mathbf{d}}^T$ |

Finally, we have to note that the approach used in the derivation of the embedded formula above is different from the one used in [8, 9, 13, 15]. By this approach of constructing embedded EPTRKN methods, there exist several embedded formulas for an EPTRKN method.

2.3. Stability Properties

Stability of (constant stepsize) EPTRKN methods was investigated by applying them to the model test equation $y''(t) = \lambda y(t)$, where $\lambda$ runs through the eigenvalues of the Jacobian matrix $\partial \mathbf{f}/\partial \mathbf{y}$ which are assumed to be negative real. It is characterized by the spectral radius $\rho(M(x))$, $x = \lambda h^2$, of the $(s+2)\times(s+2)$ *amplification matrix* $M(x)$ defined by (cf. [14, Sec. 2.2])

$$M(z) = \begin{pmatrix} xA & \mathbf{e} & \mathbf{c} \\ x^2\mathbf{b}^T A & 1 + x\mathbf{b}^T\mathbf{e} & 1 + x\mathbf{b}^T\mathbf{c} \\ x^2\mathbf{d}^T A & x\mathbf{d}^T\mathbf{e} & 1 + x\mathbf{d}^T\mathbf{c} \end{pmatrix}. \tag{2.7a}$$

The stability interval of an EPTRKN method is given as

$$(-\beta_{stab}, 0) := \{x : \rho(M(x)) \leqslant 1\}. \tag{2.7b}$$

The stability intervals of the EPTRKN methods used in our numerical codes can be found in Sec. 3.

## 3. Numerical Experiments

In this section we shall report the numerical results obtained by the sequential code ODEX2 and two our new parallel EPTRKN codes for comparing their efficiency.

### 3.1. Specifications of the Codes

ODEX2 is an extrapolation code for special second-order ODEs of the form (1.1). It uses variable order and stepsize and is implemented in the same way as the ODEX code for first-order ODEs (cf. [24, p. 294, 298]). This code is recognized as being one of the most efficient sequential integrators for nonstiff problems like (1.1) (see [24, p. 484]). In the numerical experiments, we apply ODEX2 code with standard parameter settings.

Our first code uses a variable stepsize embedded EPTRKN method based on collocation vector $\mathbf{c} = (c_1, c_2, c_3, 1)^T$ which satisfies the relations

$$\int_0^1 x^{j-1} \prod_{i=1}^{4} (x - c_i)dx = 0, \quad j = 1, 2 \tag{3.1a}$$

$$(\mathbf{b}^T + \mathbf{d}^T)\Big[\frac{\mathbf{c}^{s+2}}{s+2} - A(s+1)(\mathbf{c} - \mathbf{e})^s\Big] = 0, \tag{3.1b}$$

where (3.1a) is an orthogonality relation (cf. [24, p. 212]), and (3.1b) is introduced for minimizing the stage error coefficients (cf. [29]). The resulting method is of step point order 6 and stage order 5 (see Theorem 2.1). It has 4 as the optimal number of processors, and an embedded formula of order 3 (see Theorem 2.2). Its stability interval as defined in Sec. 2.3 is determined by numerical search techniques to be $(-0.720, 0)$. This first code is denoted by EPTRKN4.

Our second code uses a variable stepsize embedded EPTRKN method based on collocation vector $\mathbf{c} = (c_1, \ldots, c_8)^T$ which is obtained by solving the system of equations

$$\int_0^1 x^{j-1} \prod_{i=1}^{8} (x - c_i) dx = 0, \quad j = 1, 2, 3, \tag{3.2a}$$

$$c_4 = 1, c_{4+k} = 1 + c_k, \quad k = 1, 2, 3, 4. \tag{3.2b}$$

Here (3.2a) is again an orthogonality relation. The resulting method is of step point order 10 and stage order 9 (see also Theorem 2.1). It has 8 as the optimal number of processors, and an embedded formula of order 7 (see also Theorem 2.2). Its stability interval is also determined by the numerical search techniques to be $(-0.598, 0)$. This second code is denoted by EPTRKN8.

Table 1 summarizes the main characteristics of the codes: the step point order $p$, the embedded order $\hat{p}$, the optimal number of processors $np$ and the stability interval $(-\beta_{stab}, 0)$.

*Table 1.* EPTRKN codes used in the numerical experiments

| Code names | $p$ | $\hat{p}$ | $np$ | $(-\beta_{stab}, 0)$ |
|---|---|---|---|---|
| EPTRKN4 | 6 | 3 | 4 | $(-0.720, 0)$ |
| EPTRKN8 | 10 | 7 | 8 | $(-0.598, 0)$ |

Both codes EPTRKN4 and EPTRKN8 are implemented using local extrapolation and direct PIRKN methods based on the same collocation points (cf. [3]) as a starting procedure. The local error of order $\hat{p}$ denoted by $LERR$ is estimated as

$$LERR = \sqrt{\frac{1}{d} \sum_{i=1}^{d} \left[ \left( \frac{y_{n+1,i} - \hat{y}_{n+1,i}}{ATOL + RTOL|y_{n+1,i}|} \right)^2 + \left( \frac{y'_{n+1,i} - \hat{y}'_{n+1,i}}{ATOL + RTOL|y'_{n+1,i}|} \right)^2 \right]}.$$

The new stepsize $h_{n+1}$ is chosen as

$$h_{n+1} = h_n \cdot \min \left\{ 2, \max \left\{ 0.5, 0.85 \cdot LERR^{-1/(\hat{p}+1)} \right\} \right\}. \tag{3.3}$$

The constants 2 and 0.5 serve to keep the stepsize ratios $\tau_{n+1} = h_{n+1}/h_n$ to be in the interval $\left[0.5, 2\right]$.

The computations were performed on a HP-Convex X-Class Computer. The parallel codes EPTRKN4 and EPTRKN8 were implemented in sequential and parallel modes. They can be downloaded from http://www.mathematik.uni-halle.de/institute/numerik/software.

### 3.2. Numerical Comparisons

The numerical comparisons in this section are mainly made in terms of computing time for an accuracy received. However, since the parameters of the two

EPTRKN methods used in this paper are new, we would like to test the performance of these methods by comparing the number of **f**-evaluations for a given accuracy.

### Test Problems

For comparing the number of **f**-evaluations, we take two very well-known small test problems from the RKN literature:
**FEHL** - the nonlinear Fehlberg problem (cf. e.g., [16, 17, 19, 20])

$$\frac{d^2\mathbf{y}(t)}{dt^2} = \begin{pmatrix} -4t^2 & -\frac{2}{\sqrt{y_1^2(t)+y_2^2(t)}} \\ \frac{2}{\sqrt{y_1^2(t)+y_2^2(t)}} & -4t^2 \end{pmatrix} \mathbf{y}(t),$$
$$\mathbf{y}(\sqrt{\pi/2}) = (0,1)^T, \quad \mathbf{y}'(\sqrt{\pi/2}) = (-2\sqrt{\pi/2}, 0)^T \quad \sqrt{\pi/2} \leqslant t \leqslant 10,$$

with highly oscillating exact solution given by $\mathbf{y}(t) = (\cos(t^2), \sin(t^2))^T$.
**NEWT** - the two-body gravitational problem for Newton's equation of motion (see e.g., [30, p. 245], [27, 20])

$$\frac{d^2y_1(t)}{dt^2} = -\frac{y_1(t)}{\left(\sqrt{y_1^2(t)+y_2^2(t)}\right)^3}, \quad \frac{d^2y_2(t)}{dt^2} = -\frac{y_2(t)}{\left(\sqrt{y_1^2(t)+y_2^2(t)}\right)^3},$$
$$y_1(0) = 1-\varepsilon, \quad y_2(0) = 0, \quad y_1'(0) = 0, \quad y_2'(0) = \sqrt{\frac{1+\varepsilon}{1-\varepsilon}}, \quad 0 \leqslant t \leqslant 20.$$

The solution components are $y_1(t) = \cos(u(t)) - \varepsilon$, $y_2(t) = \sqrt{(1+\varepsilon)(1-\varepsilon)}\,\sin(u(t))$, where $u(t)$ is the solution of Kepler's equation $t = u(t) - \varepsilon\,\sin(u(t))$ and $\varepsilon$ denotes the eccentricity of the orbit. In this example, we set $\varepsilon = 0.9$.

For comparing the computing time, we take the following three "expensive" problems:
**PLEI** - the celestial mechanics problem from [24] which models the gravity forces between seven stars in 2D space. This modelling leads to a second-order ODE system of dimension 14. Because this system is too small, it is enlarged by a scaling factor $ns = 500$ to become the new one

$$\mathbf{e} \otimes \mathbf{y}''(t) = \mathbf{e} \otimes \mathbf{f}(t, \mathbf{y}(t)), \quad \mathbf{e} \in \mathbb{R}^{ns}.$$

**MOON** - the second celestial mechanics example which is formulated in a similar way for 101 bodies in 2D space with coordinates $x_i, y_i$ and masses $m_i$ ($i = 0, \ldots, 100$)

$$x_i'' = \gamma \sum_{j=0, j\neq i}^{100} m_j (x_j - x_i)/r_{ij}^3, \quad y_i'' = \gamma \sum_{j=0, j\neq i}^{100} m_j (y_j - y_i)/r_{ij}^3,$$

where

$$r_{ij} = ((x_i - x_j)^2 + (y_i - y_j)^2)^{1/2}, \quad i, j = 0, \ldots, 100$$
$$\gamma = 6.672, m_0 = 60, m_i = 7 \cdot 10^{-3}, \quad i = 1, \ldots, 100.$$

We integrate for $0 \leqslant t \leqslant 125$ with the initial data

$$\begin{aligned}
&x_0(0) = y_0(0) = x_0'(0) = y_0'(0) = 0\\
&x_i(0) = 30\cos(2\pi/100i) + 400, \quad x_i'(0) = 0.8\sin(2\pi/100i)\\
&y_i(0) = 30\sin(2\pi/100i), \quad y_i'(0) = -0.8\cos(2\pi/100i) + 1.
\end{aligned}$$

Here no scaling was needed because the right-hand side functions are very expensive.

**WAVE** - the semidiscretized problem for 1D hyperbolic equations (see [25]).

$$\begin{aligned}
&\frac{\partial^2 u}{\partial t^2} = gd(x)\frac{\partial^2 u}{\partial x^2} + \frac{1}{4}\lambda^2(x,u), \quad 0 \leqslant x \leqslant b, \quad 0 \leqslant t \leqslant 10,\\
&\frac{\partial u}{\partial x}(t,0) = \frac{\partial u}{\partial x}(t,b) = 0,\\
&u(0,t) = \sin\Big(\frac{\pi x}{b}\Big), \qquad \frac{\partial u}{\partial t}(0,x) = -\frac{\pi}{b}\cos\Big(\frac{\pi x}{b}\Big)
\end{aligned}$$

with

$$d(x) = 10\Big[2 + \cos\Big(\frac{2\pi x}{b}\Big)\Big], \quad \lambda = \frac{4\cdot 10^{-4} g|u|}{d(x)}, \quad g = 9.81, \quad b = 1000.$$

By using second-order central spatial discretization on a uniform grid with 40 inner points we obtain a nonstiff ODE system. In order to make this problem more expensive, we enlarge it by a scaling factor $ns = 100$.

**Results and Discussion**

The three codes ODEX2, EPTRKN4 and EPTRKN8 were applied to the above test problems with $ATOL = RTOL = 10^{-1}, 10^{-2}, \ldots, 10^{-11}, 10^{-12}$. The number of sequential **f**-evaluations (for **FEHL** and **NEWT** problems) and the computing time (for **PLEI**, **MOON** and **WAVE** problems) are plotted as a function of the global error $ERR$ at the end point of the integration interval defined by

$$ERR = \sqrt{\frac{1}{d}\sum_{i=1}^{d}\Big(\frac{y_{n+1,i} - y(t_{n+1})_i}{ATOL + RTOL|y(t_{n+1})_i|}\Big)^2}.$$

For problems **PLEI**, **MOON** and **WAVE** without exact solutions in a closed form, we use the reference solution obtained by ODEX2 using $ATOL = RTOL = 10^{-14}$.

For problems **FEHL** and **NEWT** where we compare the number of **f**-evaluations for a given accuracy, the results in Fig. 1 – 2 show that on a sequential implementation mode (symbols associated with ODEX2, EPTRKN4 and EPTRKN8) the three codes are comparable. But on a parallel implementation mode, the two parallel codes EPTRKN4 and EPTRKN8 using the optimal number of processors 4 and 8, respectively (symbols associated with EPTRKN4(4) and EPTRKN8(8)) are by far superior to ODEX2, and the code EPTRKN8 is the most efficient.

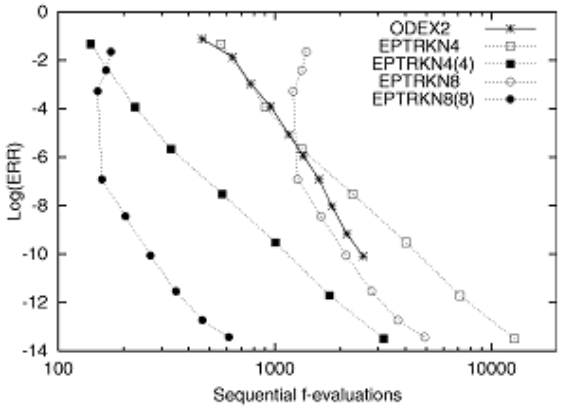


*Fig. 1.* Results for **FEHL**

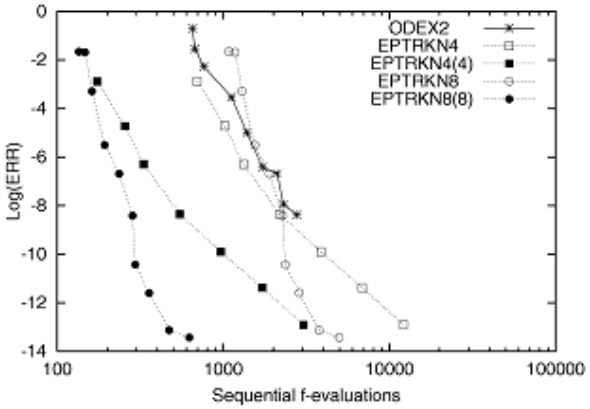


*Fig. 2.* Results for **NEWT**

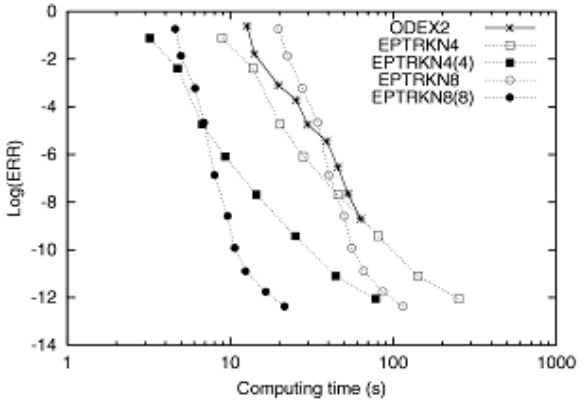


*Fig. 3.* Results for **PLEI**

For **PLEI**, **MOON** and **WAVE** problems where we compare the computing time, the results plotted in Fig. 3 – 5 show that for **PLEI** and **WAVE** problems, the two EPTRKN codes are competitive with or even more efficient than ODEX2 in the sequential implementation mode. But the parallelized EPTRKN4 and EPTRKN8 codes are again superior to ODEX2, and the results for EPTRKN4 and EPTRKN8 are almost comparable.

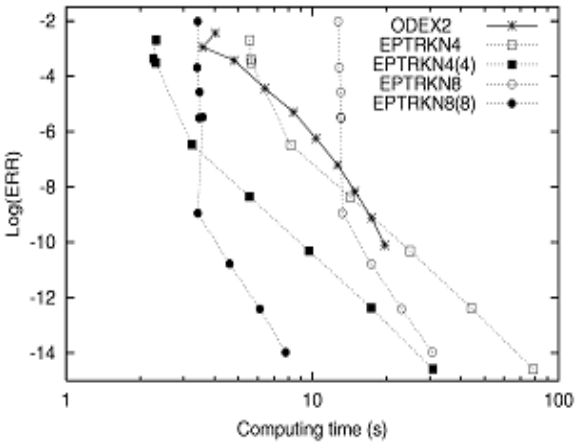


*Fig. 4.* Results for **WAVE**

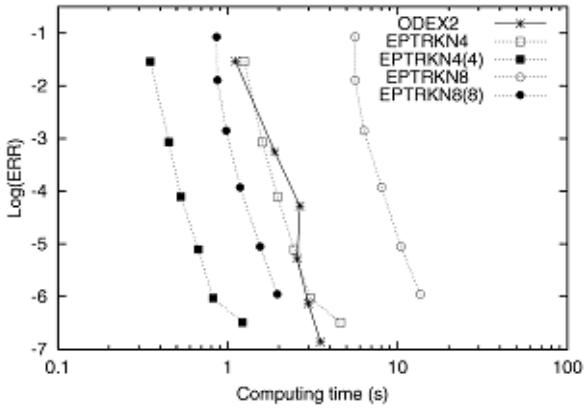


*Fig. 5.* Results for **MOON**

For **MOON**, EPTRKN4 is again competitive with ODEX2 in the sequential implementation mode. Compared with ODEX2, the parallelized EPTRKN4 and EPTRKN8 show the same efficiency as for the problems **PLEI** and **WAVE**.

The parallel speedup (cf. e.g., [13]) with $ATOL = RTOL = 10^{-8}$ as shown in Fig. 6, is very problem dependent. Using the optimal number of processors, the best speedup obtained by the codes EPTRKN4 and EPTRKN8 for the

problem **MOON** is approximately 3.3 and 5.5, respectively. We can also see from the results in Fig. 3 – 5 that for more stringent tolerances, a better speedup can be achieved.

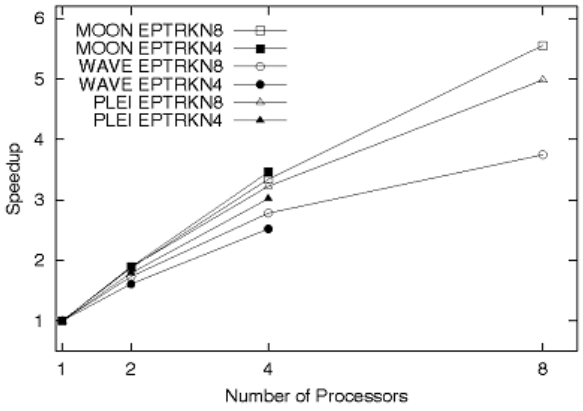


*Fig. 6.* parallel speedup

## 4. Concluding Remarks

In this paper we have considered the efficiency of a class of parallel explicit pseudo two-step RKN methods (EPTRKN methods) by comparing two new codes from this class EPTRKN4 and EPTRKN8 with the highly efficient sequential code ODEX2. By using nonstiff, expensive problems and by implementing these codes on a shared memory computer, we have shown the superiority of the new parallel codes over ODEX2. In the future we shall further improve these new parallel codes by some optimal choices of method parameters.

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