Vietnam Journal of MATHEMATICS © VAST 2006

## Generating of Minimal Unavoidable Sets

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Dedicated to Professor Do Long Van on the occasion of his 65<sup>th</sup> birthday

Received August 28, 2006 Revised October 4, 2006

Abstract. In this paper we investigate several transformations on unavoidable sets which preserve the minimality of such sets. The main result confirms that any minimal unavoidable set over an alphabet A can be obtained from A, regarded as the initial manimal unavoidable set, by finitely many applications of such transformations. As a consequence, a procedure to generate all possible minimal unavoidable sets over A is proposed. This allows in particular to generate easily counter-examples for both the Ehrenfeucht's conjecture and Haussler's one on unavoidable sets.

2000 Mathematics Subject Classification: 68R15, 68S05 Keywords: Unavoidable set, transformation, reduced, minimal, generating, conjecture.

#### 1. Introduction

In this section we recall some definitions and notations concerning with unavoidable sets and with two well-known conjectures on these sets: the *Ehrenfeucht's conjecture* and *Haussler's conjecture*. For more background we refer to [2, 5].

All the alphabets considered in this paper are supposed to be finite. Given an alphabet A, we denote by  $A^*$  the free monoid generated by A and we put  $A^+ = A^* - \{\varepsilon\}$ , where  $\varepsilon$  is the empty word. For any  $x, y \in A^*$ , we denote by  $xy^{-1}$  (resp.  $y^{-1}x$ ) the word z satisfying the condition zy = x (resp. yz = x) if it exists). Then, for any  $X, Y \subseteq A^*$  the right quotient of X by Y, denoted by  $XY^{-1}$ , is the set of all  $xy^{-1}$  with  $x \in X$  and  $y \in Y$ . The left quotient of X by Y, denoted by  $Y^{-1}X$ , is defined similarly.

Given  $X \subseteq A^*$  and  $u, w \in A^*$ . We say that w w meets u if w contains u as

a factor, i.e. w = xuy for some  $x, y \in A^*$ . We say that w avoids X if there does not exist any u in X such that w meets u.

**Definition 1.1.** Given an alphabet A and  $X \subseteq A^+$ . The set X is called an unavoidable set over A if all words in  $A^*$ , except for a finite number of them, have factors in X, or equivalently, if there exist only finitely many words avoiding X, i.e. the set  $A^* - A^*XA^*$  is finite. A set X not being an unavoidable set is called an avoidable set.

Given  $X \subseteq A^*$ ,  $u \in X$  and u' = au or u' = ua for some  $a \in A$ . Then the set  $X' = X - \{u\} \cup \{u'\}$  is called an *extention* of X at u by a on the left or on the right according as u' = au or u' = ua. Also u' is called an extention of u by a on the left or on the right according to the case.

**Definition 1.2.** Let X be an unavoidable set over an alphabet A, and |X| = n. We say that

- 1) X is extendible if there exists an extention X' of X,  $X' = X \{u\} \cup \{u'\}$ , which is still an unavoidable set.
- 2) X is n-minimal (or simply minimal) if  $\not\exists u \in X$  such that  $X \{u\}$  is still an unavoidable set.
- 3) X is n-reductive (or simply reductive) if X is n-minimal and  $\not \exists u \in X$  such that au' = u or u'a = u for some  $a \in A$ .

#### Example 1.1.

- (1)  $X = \{a^2, ab, b^2\}$  is an unavoidable set over  $A = \{a, b\}$  because  $A^* A^*XA^* = \{\varepsilon, a, b, ba\}$  which is a finite set. It is easily verified that X is a reduced unavoidable set.
- (2)  $X = \{a^3, ab, b^2\}$  is an unavoidable set over  $A = \{a, b\}$  because  $A^* A^*XA^* = \{\varepsilon, a, a^2, b, ba, ba^2\}$ . It can be verified that X is a minimal but not reduced unavoidable set.
- (3)  $X = \{a^3, ab^2, b^3\}$  is not an avoidable set because all the words of the form  $w = (ab)^n$ , the number of which is infinite, avoid X.

Remark 1.1. When replacing in X a word u by an extention u' of it, the possibility of avoiding X for any word w does not decrease, and therefore the possibility for X to be unavoidable does not increase.

Now we recall two well-known conjectures on unavoidable sets [3].

#### Ehrenfeucht's conjecture

For every unavoidable set over an alphabet A, there exists always an extension X' of X which remains an unavoidable set over A. In other words, every unavoidable set is extendible.

#### Haussler's conjecture

For any reduced unavoidable set X over an alphabet A, the maximal word-lenth of X must be smaller or equal to the cardinality |X| of X.

The following results are due to Choffrut and Culik II [1].

**Proposition 1.1.** Every unavoidable set contains a finite unavoidable set. In particular, every minimal unavoidable sets is finite.

**Theorem 1.1.** The Ehrenfeucht's conjecture is true if and only if it is true for the case of two-letter alphabets.

In [6] Rosaz has constructed a counter-example for Ehrenfeucht's conjecture. In [4], by another approach, namely by introducing and studying in deep the sets  $S_X(u,v)$  (see Definition 1.4), the authors have obtained some other counter-examples for both Ehrenfeucht's Conjecture. In this paper, we investigate several transformations on unavoidable sets which preserve the minimality of such sets. The main result confirms that any minimal unavoidable set over an alphabet A can be obtained from A – the initial minimal unavoidable set – by finitely many applications of such transformations. As a consequence, a procedure to generate all possible minimal unavoidable sets over A is proposed. This allows in particular to generate easily counter-examples for the Ehrenfeucht's and Hausler's conjectures on unavoidable sets.

**Convention:** By Theorem 1.1, from now on we may restrict ourselves to consider only the case of two-letter alphabets, namely  $A = \{a, b\}$ .

For any subset  $X \subseteq A^+$  we say that

- 1. X is a prefix set if  $\exists u \in X : u = vx \text{ with } v \in X, v \neq u \text{ and } x \in A^*$ .
- 2. X is a suffix set if  $\exists u \in X : u = xv \text{ with } v \in X, v \neq u \text{ and } x \in A^*$ .
- 3. X is an infix set if  $\exists u \in X : u = xv \text{ with } v \in X, v \neq u \text{ and } x, y \in A^*$ .
- 4. X is a bifix set if X is suffix and prefix.

#### Definition 1.3.

- 1) For u, v, w in  $A^*$ , if w = ux = yv for some  $x, y \in A^*$ ,  $xy \neq \varepsilon$ , then w is called a u-v arrow.
- 2) For any u-u arrow w, we define  $w_n = w.(w')^n$ , where  $w' = u^{-1}w$ ,  $n \ge 1$  (Figure 1).
- 3) For any u-v arrow w, we say that w avoids X if  $w \notin A^+XA^+$ .
- 4) For any u-v arrow w, we say that w is X-atomic if w avoids  $X \cup \{u, v\}$ .

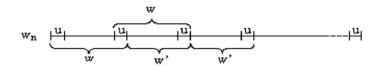


Fig. 1. Definition of u-u arrow  $w_n$  basing on u-u arrow w

**Lemma 1.1.** Let X be an infix set and  $u \in X$ . Then, the existence of a X-atomic u-u arrow implies the existence of arbitrarily long u-u arrows avoiding  $X - \{u\}$ .

*Proof.* Let W be a X-atomic u-u arrow. Let us consider a u-u arrow  $w_n = w.(w')^n$ , where  $w' = u^{-1}w$ . Because w avoids X and X is infix,  $w_n$  avoids  $X - \{u\}$ . With n large enough,  $w_n$  is the u-u arrow required.

**Lemma 1.2.** Let X be an unavoidable set. Let  $u \in X$  and u is not redundant, i.e.  $X - \{u\}$  is not an unavoidable set. Then for all words w, if w is long enough and w avoids  $X - \{u\}$  then w meets u.

*Proof.* Because u is not redundant,  $X - \{u\}$  is an avoidable set. This means that there are infinitely many words avoiding  $X - \{u\}$ . Among such words there are only finitely many words avoiding u, otherwise X is no more an unavoidable set, that contradicts the hypothesis. Thus, if w is long enough and avoids  $X - \{u\}$  then w must meet u.

**Consequence 1.1.** For each minimal unavoidable set X, there exists a natural number  $N_0$  such that for all  $w \in A^+$  with  $|w| \ge N_0$ , if  $w \notin A^*(X - \{u\})A^*$  then  $w \in A^*uA^*$  for any u in X.

**Definition 1.4.** Let X be a subset of  $A^*$ . For each pair of words u, v in X, we associate a set  $S_X(u, v)$  consisting of all X-atomic u-v arrows:

$$S_X(u,v) = uA^+ \cap A^+v - A^+XA^+$$

which we write simply S(u, v) when no confusion may arise.

**Proposition 1.2.** Let  $X \subseteq A^+$ , be an unavoidable set. X is minimal if and only if hold the following conditions:

- (i) X is an infix set.
- (ii) For all  $u \in X$ ,  $S(u, u) \neq \emptyset$ .

*Proof.* ( $\Rightarrow$ ) Suppose X is minimal, n = |X|. We will prove that the conditions (i) and (ii) must be satisfied.

- (i) If the condition (i) does not hold then there exist two words  $u, v \in X$  such that u is a proper factor of v. Then  $X \{v\}$  is also an unavoidable set, which contradicts the minimality of X. Thus (i) must hold.
- (ii) Now assume that the condition (ii) does not hold, i.e. there exists u in X such that  $S(u,u)=\emptyset$ . Because X is a minimal unavoidable set, by Consequence 1.1, there exists  $N_0$  such that for any  $w\in A^+$ , if  $|w|\geq N_0$  and w avoids  $X-\{u\}$  then w meets u. Let w be such a word. Consider the word  $w^2=w.w$ . Because w contains at least one u as factor and w avoids  $X-\{u\}$ , from  $w^2$  we can always extract a u-u arrow that is X-atomic. This means  $S(u,u)\neq\emptyset$ , a contradiction. Thus (ii) must be hold.
- ( $\Leftarrow$ ) Let X be an infix set and  $S(u, u) \neq \emptyset$  for all  $u \in X$ . Since  $S(u, u) \neq \emptyset$ , there exists a u-u arrow w which is X-atomic. By Lemma 1.1, there exists arbitrarily long u-u one which avoids  $X \{u\}$ . This means that, for any u in X,  $X \{u\}$  is not an unavoidable set. Thus X is a minimal unavoidable set.

In fact, Proposition 1.2 is a part of the following result proved in [4] in another way.

**Theorem 1.2.** [4] A set X is unavoidable set if and only if the set S(u, v) is finite for any u, v in X, and the unavoidable set X is minimal if and only if X is an infix set and S(u, u) is not empty for all u in X.

#### 2. Generating Minimal Unavoidable Sets

In this section we introduce and consider transformations on unavoidable sets which preserve the minimality, and show that every minimal unavoidable set over the alphabet  $A = \{a, b\}$  can be obtained from A, regarded as the *initial* minimal unavoidable set, by a finite number of applications of such transformations.

First we introduce a label-assigning function, denoted by Asg, which assigns to every word u in an unavoidable set X a pair  $(l_u, r_u)$  of labels called the left and the right label of u respectively.

**Definition 2.1.** Let X be a minimal unavoidable set on the alphabet  $A = \{a, b\}$ . For every word u in X,  $Asg_X(u) = (l_u, r_u)$  is defined as follows, where S stands for S(u, u).

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+ If \ suffix(S) \cap Au = \{au, bu\} \ then \ l_u = Lab.
+ If \ suffix(S) \cap Au = \{au\} \ and \ |S(u, u)| = 1 \ then \ l_u = Lia.
+ If \ suffix(S) \cap Au = \{bu\} \ and \ |S(u, u)| = 1 \ then \ l_u = Lib.
+ If \ suffix(S) \cap Au = \{au\} \ and \ |S(u, u)| > 1 \ then \ l_u = La.
+ If \ suffix(S) \cap Au = \{bu\} \ and \ |S(u, u)| > 1 \ then \ l_u = Lb.
+ If \ prefix(S) \cap uA = \{ua, ub\} \ then \ r_u = Rab.
+ If \ prefix(S) \cap uA = \{ua\} \ and \ |S(u, u)| = 1 \ then \ r_u = Ria.
+ If \ prefix(S) \cap uA = \{ua\} \ and \ |S(u, u)| > 1 \ then \ r_u = Rab.
+ If \ prefix(S) \cap uA = \{ua\} \ and \ |S(u, u)| > 1 \ then \ r_u = Rab.
+ If \ prefix(S) \cap uA = \{ua\} \ and \ |S(u, u)| > 1 \ then \ r_u = Rab.
+ If \ prefix(S) \cap uA = \{ua\} \ and \ |S(u, u)| > 1 \ then \ r_u = Rab.
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The main result of this section is the following:

**Theorem 2.1.** Let X be a minimal unavoidable set, where n = |X|. Let  $u \in X$ , and  $Asg_X(u) = (l_u, r_u)$ .

- (a) If  $l_u = Lab$ , then X can not be extended at u on the left to get a new n-minimal unavoidable set, although the set  $X' = X \{u\} \cup \{au, bu\}$  is also an (n+1)-minimal unavoidable set.
- (b) If  $l_u = Lia$ , then X can be extended consecutively on the left, beginning at u by the letter a and then by appropriate letters, to obtain infinitely many new n-minimal unavoidable sets. Similarly for the case  $l_u = Lib$ .
- (c) If  $l_u = La$ , then X can be extended consecutively on the left only finitely many times, beginning at u by the letter a and then by appropriate letters, to obtain n-minimal unavoidable sets. Similarly for the case  $l_u = Lb$ .
- (d) If  $r_u = Rab$ , then X can not be extended at u on the right to obtain a new n-minimal unavoidable set, although the set  $X' = X \{u\} \cup \{au, bu\}$  is also an (n+1)-minimal unavoidable set.

- (e) If  $r_u = Ria$ , then X can be extended consecutively on the right, beginning at u by the letter a and then by appropriate letters, to obtain infinitely many new n-minimal unavoidable sets. Similarly for the case  $r_u = Rib$ .
- (f) If  $r_u = Ra$ , then X can be extended consecutively on the right only finitely many times, beginning at u by the letter a and then by appropriate letters, to obtain new n-minimal unavoidable sets. Similarly for the case  $r_u = Rb$ .
- *Proof.* (a) A u-u arrows w is said to be of a-form if w = uy = xau with  $x \in A^*$ ,  $y \in A^*$ ,  $a \in A$ . Arrows of b-form are defined similarly. By Proposition 1.2,  $S(u, u) \neq \emptyset$  for all  $u \in X$ . Because  $l_u = Lab$ , by Definition 2.1, we have suffix $(S) \cap Au = \{au, bu\}$ . This means that S(u, u) contains u-u arrows of both a-form and b-form. We consider two possibilities:
- + Extending X at u by the letter a on the left:  $X' = X \{u\} \cup \{au\}$ . Choose in S(u,u) a u-u arrow w of b-form. Consider an arrow  $w_n$  with n large enough. Because w is X-atomic, w avoids  $X \{u\}$ . It follows from the minimality of X that X is an infix set. Hence, by Lemma 1.1, the u-u arrows  $w_n$  avoid  $X \{u\}$ . Because w is X-atomic and w is of b-form,  $w_n$  avoids au. Thus  $w_n$  avoids X', therefore X' is an unavoidable set.
- + Extending X at u by the letter b on the left  $X' = X \{u\} \cup \{bu\}$ . In a similar way we can prove that X' is not an unavoidable.

Thus, X can not be extended on the left to obtain a new n-minimal unavoidable set.

Now we prove that set  $X' = X - \{u\} \cup \{au\} \cup \{bu\}$  is a (n+1)-minimal unavoidable set. Obviously, every word long enough which avoids both au and bu also avoids u (note that the alphabet A consists of only two letters a and b). Therefore X' is also an unavoidable set. Now we prove that X' is n+1-minimal.

Since  $S_X(u,u) \neq \emptyset$  and suffix $(S) \cap Au = \{au,bu\}$ , we have  $S_{X'}(au,au) \neq \emptyset$  and  $S_{X'}(bu,bu) \neq \emptyset$ . Because X is a minimal unavoidable set, it follows by Proposition 1.2 that,  $S_X(v,v) \neq \emptyset$ , for all  $v \in X$ ,  $v \neq u$ . It is easily seen that all the words avoiding u also avoid u and u. This implies that all the words avoiding u also avoid u and u and u are u and u and u and u and u and u are u and u and u and u are u and u and u are u are u and u are u and u are u are u and u and u are u are u are u and u are u are u and u are u are u are u and u are u are u and u are u are u are u are u and u are u are u are u are u are u and u are u are u and u are u are

(b) Let  $X' = X - \{u\} \cup \{au\}$ . We shall prove that X' is n-minimal unavoidable set.

Indeed, by Consequence 1.1,  $\exists N_0 \colon \forall w \in A^+$ ,  $|w| \geq N_0$ , if w avoids  $X - \{u\}$  then w meets u. Consider an arbitrary word v avoiding  $X - \{u\}$  with  $|v| \geq 2N_0$ , we have  $v = v_1 v_2$ , where  $|v_1|, |v_2| \geq N_0$ . Since  $v, v_1, v_2$  meet u and avoid  $X - \{u\}$ , there is a factor w of v which belongs to  $S_X(u, u)$ . Because au is a suffix of w, v meets au. Thus, X' is unavoidable. Since X is minimal, there are infinitely many words v with  $|v| \geq N_0$  which avoid  $X - \{u\}$  and do not avoid v. Now, if  $|v| \geq 2N_0$  then v avoids v and v meets v. Therefore v is minimal.

Since  $S_X(u,u)$  consists of only one word w, whose suffix is au,  $S_{X'}(au,au)$  has the same form, namely  $|S_{X'}(u,u)| = 1$  and suffix  $(S_{X'}(u,u)) \cap Au = \{pu\}$  for p = a or p = b. Hence  $X'' = X' - \{au\} \cup \{pau\}$  is also n-minimal. Continuing this argument we can confirm that X can be infinitely extended on the left at u with the letter a at the first step.

The proof of (c), (d), (e), (f) are similar.

We consider now some basic transformations on any finite set  $X \subset A^*$ .

Let us note that if  $S_X(u, u) = \emptyset$  then u is redundant in X and we can delete it.

Now we consider the following transformations on an unavoidable set X.

- + Left extension by a, denoted La: transforming X into  $X' = X \{u\} \cup \{au\}$ , for some u in X with  $l_u = La$  or  $l_u = Lia$ . The transformation Lb is defined similarly.
- + Right extension by a, denoted Ra: transforming X into  $X' = X \{u\} \cup \{ua\}$ , for some u in X with  $r_u = Ra$  or  $r_u = Ria$ . The transformation Lb is defined in a similarly way.
- + Left extension by a, b, denoted Lab: transforming X into  $X' = X \{u\} \cup \{au, bu\}$ , for some u in X with  $l_u = Lab$ .
- + Right extension by a, b, denoted Rab: transforming X into  $X' = X \{u\} \cup \{ua, ub\}$ , for some u in X with  $r_u = Rab$ .
- + Left cutting, denote LC: transforming X into X', where X' is obtained from  $X \{u\} \cup \{u'\}$  with  $u' = A^{-1}u$ , by deleting all words in it which contain u' as a proper factor.
- + Right cutting, denote RC: transforming X into X', where X' is obtained from  $X \{u\} \cup \{u'\}$  with  $u' = uA^{-1}$ , by deleting all words in it which contain u' as a proper factor.
- + Deletion, denote by D: transforming X into X', where X' is obtained from X by deleting consecutively all redundant words in X.

For brevity, the transformations La, Lb, Lab, (Ra, Rb, Rab) are called commonly transformation L (transformation R, respectively).

As a direct consequence of Theorem 2.1 we obtain

Consequence 2.1. Applying transformation L and R on minimal unavoidable sets leads to minimal unavoidable sets again.

The proofs of the following lemmas are easy and therefore omitted.

- **Lemma 2.1.** Applying the transformations LC, RC and D on unavoidable sets lead to unavoidable sets again.
- **Lemma 2.2.** The initial unavoidable set  $A = \{a, b\}$  can be obtained from any finite unavoidable set by finitely many applications of the transformations LC, RC and D.
- **Lemma 2.3.** Let X and Y be minimal unavoidable sets such that  $Y = X \{u\} \cup \{u'\}$ , where u is a proper factor of u'. Then Y can be obtained from X by a finite number of applications of the transformations L and R.

The following lemma is somewhat more complex and need some verification.

**Lemma 2.4**. Let X be a minimal unavoidable set. If Y can be obtained from X by applying first some transformation RC or LC, and then some transformation

D, then X can be obtained from Y by a finite number of applications of the transformations L, R, LC, RC and D.

*Proof.* In the case |X| = |Y| the assertion is true by Lemma 2.3. So we treat only the case |X| > |Y|. By duality, it suffices to check only the case applying RC on some word x in X. Without loss of generality we may assume x = ua. When the first transformation is RC, by applying RC we get  $u \in Y$ . By using lemmas above, the following facts can be verified step by step:

- + The right label of u in Y is Rab
- + Each element deleted from X to get Y must be appeared in some ub-ub arrow avoiding Y.
- + Each u-u arrow avoiding  $Y \{u\}$  has to meet some deleted element in Z = X Y.
- + For any x in Z, each long enough x-x arrow avoiding  $Y \cup \{ua\}$  has to meet ub.
- + By applying the extension Rab on Y at u in Y, we get X',  $X' = Y \{u\} \cup \{ua, ub\}$ , and then by other transformations R and L we can obtain a new minimal unavoidable set Y' so that every deleted element x in Z is a factor of some ub-ub arrow in Y',  $|Y'| \ge |X|$ .
- + Y' can be also obtained from X by applying some transformations R and L.
- + Applying transformations RC, LC, D on Y' first at ub-ub arrows and then at factors of them we can get X.

Now the following main result of this section can be easily proved basing on the above mentioned lemmas.

**Theorem 2.2.** Every minimal unavoidable set over the alphabet  $A = \{a, b\}$  can be obtained from the initial minimal unavoidable set A by taking finitely many transformations R, L, LC, RC and D.

It is not difficult to see that the result remains valid for any finite alphabet.

Consequence 2.2. Every minimal unavoidable set over any finite alphabet A can be obtained from the initial minimal unavoidable set A by applying finitely many transformations R, L, LC, RC and D.

# 3. Generating of Counter-examples for Ehrenfeucht's and Haussler's Conjectures

Now we present a computer computer program generating finite minimal unavoidable set starting from the alphabet A. This allows to find out easily counterexamples for Ehrenfeucht's and Haussler's conjectures in one or two hours. The correctness of the procedure can be verified directly basing on Theorem 2.2.

#### **Procedure 3.1.** (interactive mode)

1. Input any minimal unavoidable set, say  $X_0$ ;

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2. X = X_0;
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- 3. Calculating  $Asg_X(u)$  for all u in X;
- 4. exitOK=0;
- 5. While exitOK=0 do
- 6. Begin
  - a) SelectOne(Transformation);
  - b) Case Transformation of

 $L_k$ : taking one left extension step of type L on  $x_k$  in X;

 $R_k$ : taking one right extension step of type R on  $x_k$  in X;

 $LC_k$ : cutting one letter on the left of  $x_k$  in X by an action of type  $LC_k$ ;

 $RC_k$ : cutting one letter on the right of  $x_k$  in X by an action of type  $RC_k$ ;

 $D_k$ : deleting  $x_k$  in X; { whenever  $x_k$  is redundant };

Write: Save list X to a file;

Copy: Copy list X to a buffer to re-use later if needed;

Exit: exitOK=1;

EndCase;

- 7. EndWhile;
- 8. EndProc;

We exhibit below some counter-examples for Ehrenfeucht's and Haussler's conjectures, which have been obtained with the aid of Precedure 3.1.

Example 3.1. Counter-examples for Haussler's conjecture:

$$X_n = \{aa, bbb, (ab)^n a, (ab)^n b, bbabb, bba(ba)bb, ..., bba(ba)^{n-2}bb\}, n = 2, 3, ...$$
  
Generating  $X_2$  is presented in Table 3.1.

Remark that the maximal wordlength of  $X_n$  is 2n+1 which is much larger than n=|X|. In [4] it is showed that there are infinitely many reduced unavoidable sets X whose maximal wordlength is of the size  $O(c^n)$ , where c>1 is a constant.

Table 3.1. Generating counter-example for Haussler Conjecture from  $A = \{a, b\}$ 

Step	Left label	Right label	Unavoidable sets	Next operation
1	Lia Lib	Ria Rib	a b	R1
2	Lia Lab	Ria Rab	aa b	R2
3	Lia Lia Lib	Ria Rib Rib	aa ba bb	R3
4	Lia Lab Lib	Ria Rb Rib	aa ba bbb	R2

Step	Left label	Right label	Unavoidable sets	Next operation
5	Lab Lab Lib	Rab Rab Rib	aa bab bbb	L2
6	Lab Lib Lia Lib	Rab Ria Rib Rib	aa abab bbab bbb	R2
7	Lab Lib La Lib	Rab Rib Rab Rib	aa ababa bbab bbb	R2
8	Lab Lib La Lib	Rab Ria Rab Rib	aa ababab bbab bbb	R3
9	Lab Lib Lia Lia Lib	Rab Ria Rib Ria Rib	aa ababab bbaba bbabb bbb	R3
10	Lab Lib Lia Lia Lib	Rab Ria Rib Ria Rib	aa ababab bbabab bbabb bbb	R3
11	Lab Lb Lia Lia Lib	Rab Rab Ria Ria Rib	aa ababab bbababb bbabb bbb	R2
12	Lab Lib Lib Lia Lia Lib	Rab Rib Ria Ria Ria Rib	aa abababa abababb bbababb bbabb	

Example 3.2. Counter-examples for Ehrenfeucht's conjecture:

a)  $X = \{a^3, b^4, ab^3ab, ab^2ab, abab, b^2a^2b^2, ba^2ba^2b\}$ 

(see Table 3.2). This counter-examples has been obtained first by Rozas [6].

- b)  $X = \{a^4, b^4, ba^3b, ba^2ba^2b, baba, b^2a^2b^2, bab^2a, bab^3a\},$
- c)  $X = \{a^3, b^4, a^2ba^2b, baba, b^2a^2b^2, bab^2a, bab^3a\}...$

Table 3.2. Generating counter-example for Ehrenfeut conjecture from  $\{a,b\}$ 

Step	Left label	Right label	Unavoidable sets	Next operation
1	Lia Lib	Ria Rib	a b	R1
2	Lia Lab	Ria Rab	aa b	L2
3	Lia Lab Lib	Ria Ra Rib	aaa ab bb	R3
4	Lia Lab Lib	Ria Rab Rib	aaa ab bbb	R3
5	Lia Lab Lib	Ria Rab Rib	aaa ab bbbb	R2
6	Lia Lab Lab Lib	Ria Rab Rab Rib	aaa aba abb bbbb	R2
7	Lia Lia Lib Lab Lib	Ria Rib Ria Rab Rib	aaa abaa abab abb bbbb	R4
8	Lia Lia Lib Lab Lab Lib	Ria Rib Ria Rab Ra Rib	aaa abaa abab abba abbb bbbb	R5
9	Lia Lia Lib Lab Lab Lab	Ria Rib Ria Rab Rab Rab	aaa abaa abab abba abbba bbbb	R4

Step	Left	Right	Unavoidable	Next
	label	label	sets	operation
10	Lia Lia Lib Lia Lib Lab Lab	Ria Rib Ria Rib Rib Rab Rab	aaa abaa abab abbaa abbab abbba bbbb	R4
11	Lab Lia Lib Lia Lib Lab Lab	Rab Rib Ria Rib Rib Rab Rab	aaa abaa abab abbaab abbab abbba bbbb	R2
12	Lab Lia Lib Lia Lib Lab Lab	Rab Ria Ria Rib Rib Rab Rab	aaa abaab abab abbaab abbab abbba bbbb	L2
13	Lab Lib Lab Lia Lab Lab Lab	Rab Ria Ra Rib Rab Rab Rab	aaa aabaab abab abbaab abbab abbba bbbb	R6
14	Lab Lia Lab Lia Lab Lab Lab	Rab Rib Ra Rib Rab Rab Rab	aaa baabaab abab abbaab abbab abbba bbbb	LC4
15	Lab Lia Lab Lia Lab Lab	Rab Ria Ra Rib Rab Rb	aaa baabaab abab bbaab abbab abbba bbbb	R6

label label sets operation    Comparison   Comparison	G)	T C:	D: 14	TT • 1 1 1	NT /
16 Lia Ria aaa R2 Lab Rab Abbab Abbab Lab Rab Abbab Abbab Lab Rab Abbab Abbab Lab Rab Abbab Ab	Step	Left	Right	Unavoidable	Next
Lib Rab Ra abab Rab Lab Rab Babab Babab Rab Babab Babab Rab Babbab Lab Rab Babab Babbab Ba		label	label	sets	operation
Lab Rab bbaa abbab Lab Rab Rab abbab Lab Rab Bbaa Bbab Bbaa Abbab Lab Rab Bbaa Abbab Lab Rab Bbaa Bbab Bbaa Bbab Bbaa Bbab Bbab	16	Lia	Ria	aaa	
Lab Rab abbab Abbab Lab Rab Bbaa Abbab Lab Rab Bbaa Abbab Bbaa Abbab Bbaa Abbab Bbaa Bbab Bbabb Bbab Bbab Bbabb Bbabb Bbabb Bbabb Bbabb Bbabb Bbabb Bbabb Bbabbab B		Lib	$\operatorname{Rib}$	aabaa	R2
Lab Rab abbab abbab Lab Rab Babab Abab Babab Babbab		Lab	Ra	abab	
Lab Rab bbbb		Lab	$\operatorname{Rb}$	bbaa	
LabRabbbbb17LabRabaaaLibRiaaabaabL2LabRaababLabRabbbaaLabRababbabLabRababbabLabRababbabLabRabaaaLiaRiabaabaabLabRaababLabRaabbabLabRababbabLabRababbabLabRababababLabRababababLabRaabababLabRaabababLabRababbabLabRababbabLabRababababLabRabbaabaabLabRabbaababLabRabbaababLabRabababLabRabababLabRabababLabRababbabLabRababbabLabRababbabLabRababbab		Lab	Rab	abbab	
17 Lab Ria aabaab L2 Lib Ria aabaab L2 Lab Ra abab Lab Rab bbaa Lab Rab abbab Lab Rab abbab Lab Rab abbab Lab Rab ababab Lab Rab abbab Lab Rab ababab Lab Rab ababab Lab Rab abbab Lab Rab ababab Lab Rab abbab Lab Rab abbab		Lab	Rab	abbbab	
Lib Ria aabaab Lab Ra abaab Lab Rab abbab Lab Rab abbab Baab Babab Babbab Babb		Lab	Rab	bbbb	
Lab Rab Bab Bab Bab Bab Bab Bab Bab Bab Bab B	17	Lab	Rab	aaa	
Lab Rab abbab Lab Rab Bbaa abbab Lab Rab Bab Bbaa Bbab Rab Bbab Bbaa Babab Babbab Babba		Lib	Ria	aabaab	L2
Lab Rab abbab Lab Rab bbbb  18 Lab Rab ababab Lab Rab abab Lab Baabaab Lab Baabaab Lab Baabaab Lab Baabaab Lab Baabaab Lab Bab abbab Lab Bab Baabaab Lab Bab Baabaab Lab Bab Baabaab Lab Bab Babbab		Lab	Ra	abab	
Lab Rab bbbb  18 Lab Rab bbbb  18 Lab Rab aaa Baabaab Lab Ra abab Bbab Bbab Bbab Baabaab Baabab Baab Baa			$\operatorname{Rb}$	bbaa	
LabRabbbbb18LabRabaaaLiaRiabaabaabaLabRaababLabRabbbaaR4LabRababbabLabRababbabLabRabaaaLiaRiabaabaabLabRaababLabRababbabLabRababbabLabRababbbab20LabRabbaabaababLabRabababLabRababaabLabRababaabLabRababaabLabRababbabLabRababbabLabRababbabLabRababbabLabRababbabLabRababbab					
Lab Rab Bababab R4  Lia Ria baabaab Babbab Bab Bab Babab Bab Bab B					
Lia Ria baabaab Lab Ra abab Lab Rb bbaa R4  Lab Rab abbab Lab Rab abbbab Lab Rab bbbb  19 Lab Ra abab Lab Ra abab Lab Ra abab Lab Ra abab Lab Rab abbab Lab Rab abbab Lab Rab abbbb  20 Lab Rab Rab abab Lab Rab baabaab Lab Rab baabaab Lab Rab abab Lab Rab abbab Lab Rab abbab		Lab	Rab	bbbb	
Lab Ra abab Bbaa R4  Lab Rab abbab Bbaa R4  Lab Rab abbab Bbbb  19 Lab Rab ababab Bbbb  19 Lab Rab ababab Bbaab Bbaab Bbaab Bbab Bba	18	Lab	Rab	aaa	
Lab Rab abbab R4  Lab Rab abbab Rab abbbab  Lab Rab abbbab  Lab Rab ababaab  Lab Ra abab  Lab Rab abbab R4  Lab Rab abbab  Lab Rab abbab  Lab Rab abbab  Lab Rab ababab  Lab Rab ababab  Lab Rab ababab  Lab Rab ababab  Lab Rab abab  Lab Rab abbab  Lab Rab abbab  Lab Rab abbab  Lab Rab abbab		Lia	Ria	baabaab	
Lab Rab abbab Lab Rab bbbb  19 Lab Rab baabaab Lia Ria baabaab Lab Ra abab Lab Rab abbab Lab Rab abbab Lab Rab abbab Lab Rab abbab Lab Rab baabaab Lab Rab baabaab Lab Rab baabab Lab Rab baabaab Lab Rab abab Lab Rab abbab Lab Rab abbab		Lab	Ra	abab	
Lab Rab bbbbb  19 Lab Rab bababab bbbb  19 Lab Rab aaa baabaab Lab Rab bbaab R4  Lab Rab abbab Bbaab Babab Bbabb Bbabb Bbbbb  20 Lab Rab abab Babab Babab Babab Babbab Babb Bbaab Bbabb Bbbbb Bbbbb Bbbbb Babbbbb Bbbbb Babbbbbb Babbbbbbb Babbbbbbbb		Lab	$\operatorname{Rb}$	bbaa	R4
Lab Rab bbbb  19 Lab Rab aaa Lia Ria baabaab Lab Ra abab Lab Rab bbaab R4 Lab Rab abbab Lab Rab abbbbb  20 Lab Rab Bab aaa Lab Rab baabaab Lab Rab baabaab Lab Rab abab Lab Rab abbab Lab Rab abbab		Lab	Rab	abbab	
19 Lab Rab aaa Lia Ria baabaab Lab Ra abab Lab Rab bbaab R4 Lab Rab abbab Lab Rab abbbab Lab Rab baabaab Lab Rab baabaab Lab Rab baabaab Lab Rab abab Lab Rab abbab Lab Rab abbab Lab Rab abbab		Lab	Rab	abbbab	
Lia Ria baabaab Lab Ra abab Lab Rb bbaab R4 Lab Rab abbab Lab Rab abbab Lab Rab bbabb  20 Lab Rab baabaab Lab Rab baabaab Lab Rab abab Lab Rab abab Lab Rab abab Lab Rab abab Lab Rab abbab Lab Rab abbab Lab Rab abbab		Lab	Rab	bbbb	
Lab Ra abab Lab Rb bbaab R4 Lab Rab abbab Lab Rab abbab Lab Rab bbbb  20 Lab Rab aaa Lab Rab baabaab Lab Rab abab Lab Rab abab Lab Rab abab Lab Rab abbab Lab Rab abbab Lab Rab abbab	19	Lab	Rab	aaa	
Lab Rb bbaab R4  Lab Rab abbab Lab Rab abbbab Lab Rab bbbb  20 Lab Rab baabaab Lab Rab baabaab Lab Rab abab Lab Rab abab Lab Rab abab Lab Rab abbab Lab Rab abbab Lab Rab abbab		Lia	Ria	baabaab	
Lab Rab abbab Lab Rab abbbab Lab Rab bbbb  20 Lab Rab baabaab Lab Rab baabaab Lab Rab abab Lab Rab abab Lab Rab abbab Lab Rab abbab Lab Rab abbab		Lab	Ra	abab	
Lab Rab abbbab Lab Rab aaa  Lab Rab baabaab  Lab Rab abab  Lab Rab abab  Lab Rab abbab  Lab Rab abbab  Lab Rab abbab  Lab Rab abbab		Lab	Rb	bbaab	R4
Lab Rab bbbb  20 Lab Rab aaa Lab Rab baabaab Lab Rab abab Lab Rab bbaabb Lab Rab abbab Lab Rab abbab Lab Rab abbab					
20 Lab Rab aaa Lab Rab baabaab Lab Rab abab Lab Rab bbaabb Lab Rab abbab Lab Rab abbab					
Lab Rab baabaab Lab Rab abab Lab Rab bbaabb Lab Rab abbab Lab Rab abbbab		Lab	Rab	bbbb	
Lab Rab abab Lab Rab bbaabb Lab Rab abbab Lab Rab abbbab	20				
Lab Rab bbaabb Lab Rab abbab Lab Rab abbbab					
Lab Rab abbab Lab Rab abbbab					
Lab Rab abbbab					
Lah   Rah   bbbb					
Lau Itau Duuu		Lab	Rab	bbbb	

Acknowledgements. The authors would like to thank the referee for many helpful commands which help us to improve the presentation of the paper more better than before.

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